

## Sluicing as anaphora to a scope remnant

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Synopsis: I argue that sluicing is anaphora to a continuation, that is, to a constituent missing a piece. When a DP takes scope over a clause, it creates the right kind of antecedent. The prediction is that sluicing is sensitive to scope islands, but not to overt-movement islands.



Richard Montague



Robert May

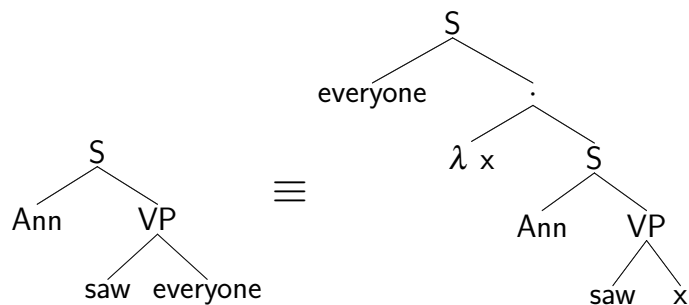
Today's question: *How to incorporate QR into a genuine logic?*

## Quantifier Raising: a logical inference?

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- **Montague** 1973: Quantifying In: (2661 citations)
- **May** 1978,1985: Quantifier Raising (QR): (2866 citations)

Montague  $\downarrow$   $\frac{\text{everyone}(\lambda x. \text{Ann saw } x) \vdash S}{\text{Ann saw everyone} \vdash S}$   $\uparrow$  May



## Lambek's substructural logic NL for natural language

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Without Exchange, ' $\rightarrow$ ' splits into ' $\backslash$ ' and ' $/$ '

- **Formulas:**  $\mathcal{F} = \text{DP} \mid S \mid \mathcal{F} \backslash \mathcal{F} \mid \mathcal{F} / \mathcal{F}$
- **Structures:**  $\mathcal{S} = \mathcal{F} \mid \mathcal{S} \cdot \mathcal{S}$
- **Sequents:**  $\mathcal{S} \vdash \mathcal{F}$
- **Logical rules:**

$$\frac{\Gamma \vdash A \quad \Sigma [B] \vdash C}{\Sigma [\Gamma \cdot A \backslash B] \vdash C} \backslash L$$

$$\frac{A \cdot \Gamma \vdash B}{\Gamma \vdash A \backslash B} \backslash R$$

$$\frac{\Gamma \vdash A \quad \Sigma [B] \vdash C}{\Sigma [B / A \cdot \Gamma] \vdash C} / L$$

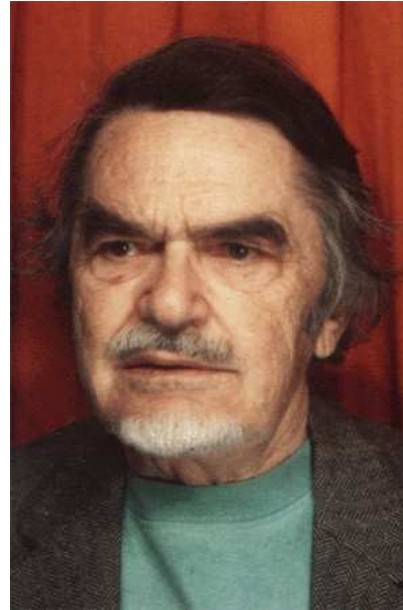
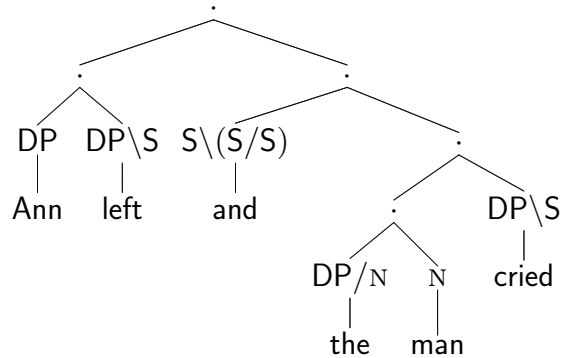
$$\frac{\Gamma \cdot A \vdash B}{\Gamma \vdash B / A} / R$$

Structural rules: none! (Cut baked in)

## How context notation works in inference rules

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- Capital Greek letters ( $\Delta, \Gamma, \Sigma$ ) stand for complete structures
- ' $\Sigma[\Delta]$ '  $\equiv$   $\Sigma$  containing a distinguished instance of  $\Delta$
- ' $\Sigma[\Gamma \cdot A \backslash B]$ ' matches the structure below in two ways:
  - $[\text{Ann} \cdot \text{DP} \backslash \text{S}] \cdot (\text{and} \ ((\text{the} \cdot \text{man}) \cdot \text{cried}))$
  - $(\text{Ann} \cdot \text{left}) \cdot (\text{and} \cdot [(\text{the} \cdot \text{man}) \cdot \text{DP} \backslash \text{S}])$



Joachim Lambek

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## An example derivation of Ann saw Bill

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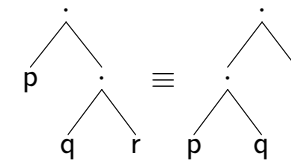
## Adding a structural rule for QR

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Associativity:  $p \cdot (q \cdot r) \equiv (p \cdot q) \cdot r$

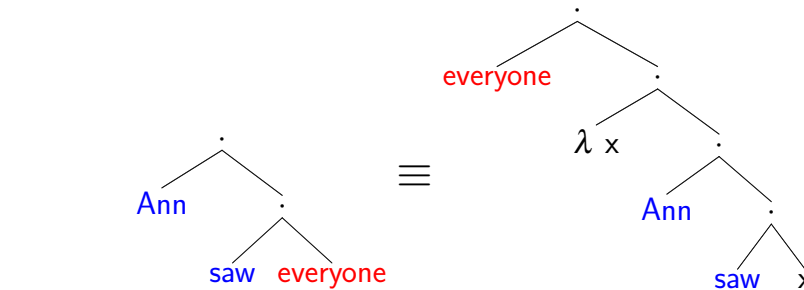
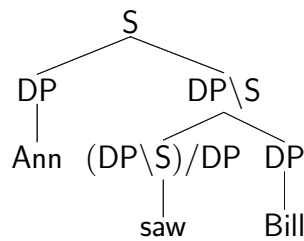
(1)

$$\frac{\frac{\frac{\text{DP} \vdash \text{DP} \quad \text{S} \vdash \text{S}}{\text{DP} \cdot \text{DP} \backslash \text{S} \vdash \text{S}} \backslash L}{\text{DP} \cdot ((\text{DP} \backslash \text{S}) / \text{DP} \cdot \text{DP}) \vdash \text{S}} / L}{\text{Ann} \cdot (\text{saw} \cdot \text{Bill}) \vdash \text{S}} \text{LEX}$$



Quantifier Raising:  $\Sigma[\Delta] \equiv \Delta \cdot \lambda x \Sigma[x]$

(2)



(3)

- Curry-Howard: L rules correspond to function application
- saw(bill)(ann)**

- **Variables:**  $\mathcal{V} = x \mid y \mid z \mid \dots$
- **Formulas:**  $\mathcal{F} = DP \mid S \mid \mathcal{F} \setminus \mathcal{F} \mid \mathcal{F} / \mathcal{F}$
- **Structures:**  $\mathcal{S} = \mathcal{F} \mid \mathcal{S} \cdot \mathcal{S} \mid \mathcal{V} \mid \lambda \mathcal{V} \mathcal{S}$
- **Sequents:**  $\mathcal{S} \vdash \mathcal{F}$
- **Logical rules:**

$$\frac{\Gamma \vdash A \quad \Sigma[B] \vdash C}{\Sigma[\Gamma \cdot A \setminus B] \vdash C} \setminus L \qquad \frac{A \cdot \Gamma \vdash B}{\Gamma \vdash A \setminus B} \setminus R$$

$$\frac{\Gamma \vdash A \quad \Sigma[B] \vdash C}{\Sigma[B / A \cdot \Gamma] \vdash C} /L \qquad \frac{\Gamma \cdot A \vdash B}{\Gamma \vdash B / A} /R$$

- **Structural rule:**  $\Sigma[\Delta] \equiv_{QR} \Delta \cdot \lambda x \Sigma[x]$

Linear: !1 var per lambda; x chosen fresh

Works great!

$$\frac{\vdots \quad \text{Ann} \cdot (\text{saw} \cdot DP) \vdash S}{DP \cdot \lambda x (\text{Ann} \cdot (\text{saw} \cdot x)) \vdash S} \text{QR}$$

$$\frac{\lambda x (\text{Ann} \cdot (\text{saw} \cdot x)) \vdash DP \setminus S \quad S \vdash S}{\lambda x (\text{Ann} \cdot (\text{saw} \cdot x)) \vdash DP \setminus S} \setminus R$$

$$\frac{S / (DP \setminus S) \cdot \lambda x (\text{Ann} \cdot (\text{saw} \cdot x)) \vdash S}{\text{everyone} \cdot \lambda x (\text{Ann} \cdot (\text{saw} \cdot x)) \vdash S} /L$$

$$\frac{\text{everyone} \cdot \lambda x (\text{Ann} \cdot (\text{saw} \cdot x)) \vdash S}{\text{Ann} \cdot (\text{saw} \cdot \text{everyone}) \vdash S} \text{QR}$$

...including the Curry-Howard labeling for the semantics:

$$\frac{\vdots \quad \text{ann} \cdot (\text{saw} \cdot y) \vdash \text{saw} y \text{ ann}}{y \circ \lambda x (\text{ann} \cdot (\text{saw} \cdot x)) \vdash \text{saw} y \text{ ann}} \text{QR}$$

$$\frac{\lambda x (\text{ann} \cdot (\text{saw} \cdot x)) \vdash \lambda y. \text{saw} y \text{ ann} \quad p \vdash p}{\lambda x (\text{ann} \cdot (\text{saw} \cdot x)) \vdash \lambda y. \text{saw} y \text{ ann}} \setminus R$$

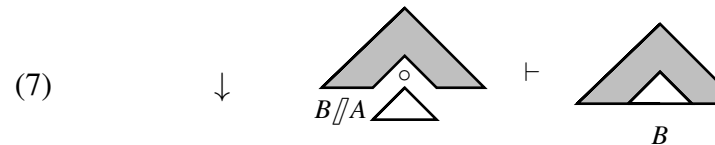
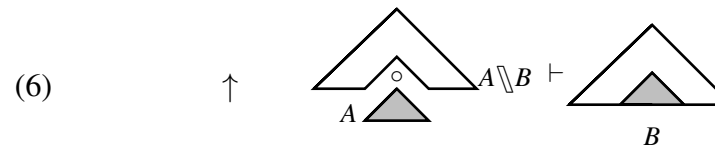
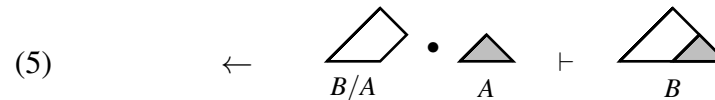
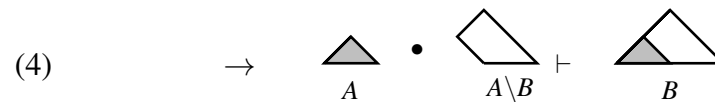
$$\frac{Q \circ \lambda x (\text{ann} \cdot (\text{saw} \cdot x)) \vdash Q(\lambda y. \text{saw} y \text{ ann})}{\text{everyone} \circ \lambda x (\text{ann} \cdot (\text{saw} \cdot x)) \vdash \text{everyone}(\lambda y. \text{saw} y \text{ ann})} /L$$

$$\frac{\text{everyone} \circ \lambda x (\text{ann} \cdot (\text{saw} \cdot x)) \vdash \text{everyone}(\lambda y. \text{saw} y \text{ ann})}{\text{ann} \cdot (\text{saw} \cdot \text{everyone}) \vdash \text{everyone}(\lambda y. \text{saw} y \text{ ann})} \text{LEX QR}$$



Michael Moortgat

Two modes of syntactic combination



Compare with tangram diagrams in Moortgat 1996b

Parasitic scope: sentence-internal same

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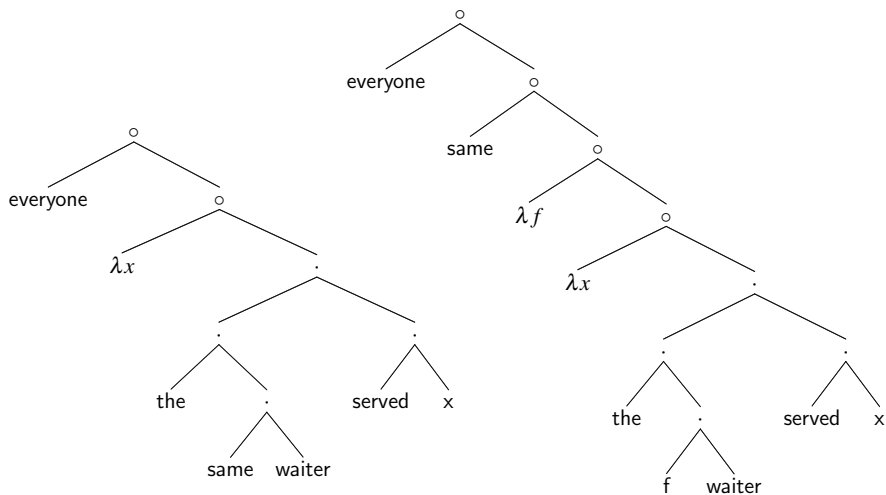
- (8) a. The same waiter served everyone. [Stump, Heim]
- b. There is a (unique) waiter  $x$  such that  $x$  served everyone.

$$\begin{array}{c}
 \vdots \\
 \frac{(the \cdot (A \cdot waiter)) \cdot (served \cdot DP) \vdash S}{DP \circ \lambda x((the \cdot (A \cdot waiter)) \cdot (served \cdot x)) \vdash S} \lambda \\
 \frac{\lambda x((the \cdot (A \cdot waiter)) \cdot (served \cdot x)) \vdash DP \backslash S}{A \circ \lambda y \lambda x((the \cdot (y \cdot waiter)) \cdot (served \cdot x)) \vdash DP \backslash S} \backslash R \\
 \frac{\lambda y \lambda x((the \cdot (y \cdot waiter)) \cdot (served \cdot x)) \vdash A \backslash (DP \backslash S)}{(DP \backslash S) // (A \backslash (DP \backslash S)) \circ \lambda y \lambda x((the \cdot (y \cdot waiter)) \cdot (served \cdot x)) \vdash DP \backslash S} \backslash R \\
 \frac{DP \backslash S \vdash DP \backslash S}{(DP \backslash S) // (A \backslash (DP \backslash S)) \circ \lambda y \lambda x((the \cdot (y \cdot waiter)) \cdot (served \cdot x)) \vdash DP \backslash S} //L \\
 \frac{same \circ \lambda y \lambda x((the \cdot (y \cdot waiter)) \cdot (served \cdot x)) \vdash DP \backslash S}{S // (DP \backslash S) \circ (same \circ \lambda y \lambda x((the \cdot (y \cdot waiter)) \cdot (served \cdot x)) \vdash S)} LEX \\
 \frac{S // (DP \backslash S) \circ (same \circ \lambda y \lambda x((the \cdot (y \cdot waiter)) \cdot (served \cdot x)) \vdash S)}{everyone \circ (same \circ \lambda y \lambda x((the \cdot (y \cdot waiter)) \cdot (served \cdot x)) \vdash S)} LEX \\
 \frac{everyone \circ (same \circ \lambda y \lambda x((the \cdot (y \cdot waiter)) \cdot (served \cdot x)) \vdash S)}{everyone \circ \lambda x((the \cdot (same \cdot waiter)) \cdot (served \cdot everyone)) \vdash S} \lambda \\
 \frac{everyone \circ \lambda x((the \cdot (same \cdot waiter)) \cdot (served \cdot everyone)) \vdash S}{(the \cdot (same \cdot waiter)) \cdot (served \cdot everyone) \vdash S} \lambda
 \end{array}$$

Details in Barker 2007; not derivable in MM96

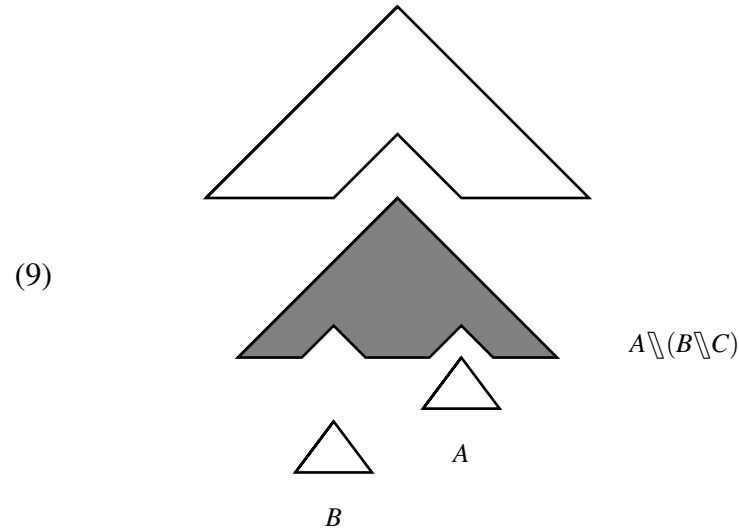
Parasitic scope in tree format

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Parasitic scope in schematic format

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Grey constituent ~ string with two points of discontinuity

Other phenomena with a parasitic scope analysis

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- (10) a. Anaphora: Morrill, Fadda & Valentín 2011
- b. *he*:  $(DP \backslash S) // (DP \backslash (DPS))$
- c. Everyone thinks he is smart.
- d.  $everyone \circ (he \circ \lambda y \lambda x(x \cdot (thinks \cdot (y \cdot (is \cdot smart)))))) \vdash S$
- (11) a. *Average*: Kennedy and Stanley 2009
- b. The average American has 2.3 kids.
- c.  $2.3 \circ (avg \circ \lambda f \lambda n((the \cdot (f \cdot Am'n)) \cdot (has \cdot (n \cdot kids))))$
- (12) a. Fancy coordination: Kubota & Levine (various papers)
- b. I said the same thing to Terry on Mon and to Kim on Tue.
- c.  $\neq$  I said the same thing to Terry on Monday and I said the same thing to Kim on Tuesday.
- (13) a. Remnant comparatives: Pollard and Smith 2013
- b. Ann owes Bill more than Clara.

Kubota and Levine's workshop in week 2!

- (14) a. Solomon 2009  
 b. Ann and Bill know [some of the **same** people].  
 c. There is a set of people  $X$  such that Ann knows some of  $X$  and Bill knows some of  $X$ .  
 d. No guarantee that Ann and Bill know anyone in common!  
 e. Solomon:  $same:((DP \setminus S) // (DP \setminus (DP \setminus S))) // (A \setminus DP)$

$$(15) \frac{\frac{\text{they} \circ ((\text{same} \circ \lambda x(\text{some} \cdot (\text{of} \cdot (\text{the} \cdot (x \cdot \text{people})))))) \circ \lambda zy(y \cdot (\text{know} \cdot z))) \vdash S}{\text{they} \circ \lambda y(y \cdot (\text{know} \cdot (\text{same} \circ \lambda x(\text{some} \cdot (\text{of} \cdot (\text{the} \cdot (x \cdot \text{people})))))) \vdash S} \lambda}{\frac{\text{they} \cdot (\text{know} \cdot (\text{same} \circ \lambda x(\text{some} \cdot (\text{of} \cdot (\text{the} \cdot (x \cdot \text{people})))))) \vdash S}{\text{they} \cdot (\text{know} \cdot (\text{some} \cdot (\text{of} \cdot (\text{the} \cdot (\text{same} \cdot \text{people})))) \vdash S} \lambda} \lambda$$

lancet liver fluke (*Dicrocoelium dendriticum*)

Sluicing as anaphora to an anti-constituent

- (1) Someone left, but I don't know [who \_].  
 (2) [Someone<sub>INNER ANTECEDENT</sub> left]<sub>OUTER ANTECEDENT</sub>,  
 but I don't know [who<sub>WH</sub> SLUICEGAP]<sub>SLUICE</sub>.

$$\begin{aligned} \text{sluice} &= \text{wh-phrase} + (\text{antecedent-clause} - \text{inner-antecedent}) \\ &= \text{who} + ([\text{someone left}] - \text{someone}) \\ &= \text{who} + [\_ \text{left}] \end{aligned}$$

- The outer antecedent with the inner antecedent removed
- The remnant of the outer antecedent after the inner antecedent has taken scope (i.e., a nuclear scope)
- The complement of the inner antecedent with respect to the outer antecedent, i.e., **an anti-constituent**
- The delimited **continuation** of the inner antecedent wrt to the outer antecedent

Three comparison analyses: structured silence?

Some analyses of sluicing assume that the sluice ellipsis site contains a silent object that has internal structure:

- **LF copying:** Chung, Ladusaw and McCloskey 1995
  - Re-use (“recycle”) the Logical Form of the antecedent
  - Builds silent structure inside sluicegap
- **PF Deletion:** Merchant 2001
  - Build any IP you want to. Move the WH out; delete the remainder if there is a certain kind of semantic equivalence with the antecedent

Other analyses propose that sluicing is a kind of anaphora:

- **Anaphora:** Jäger 2005
  - Antecedent: clause containing an indefinite
  - No internal structure to silence

Three puzzles to use for comparing analyses

**Case matching:** the case of the WH element in the sluice

- must match the case of the inner antecedent.
- (4) Er will jemandem schmeicheln, aber sie wissen nicht, { \*wen / wem }.  
 he wants someone.DAT flatter but they know not { who.ACC / who.DAT }  
 ‘He wants to flatter someone, but they don't know who.’
- (5) Er will jemanden loben, aber sie wissen nicht, { wen / \*wem }.  
 he wants someone.ACC praise but they know not { who.ACC / who.DAT }  
 ‘He wants to praise someone, but they don't know who.’

**Island insensitivity:** the inner antecedent can be embedded within an island for WH-movement.

- (6) He wants a detailed list, but I don't know how detailed [he wants a \_\_\_ list] (\* if pronounced)
- (7) Bo talked to the people who discovered something, but we don't know what [Bo talked to the people who discovered \_\_\_].

**Sprouting:** sometimes there is no overt inner antecedent

- (10) John left, but I don't know when.

## Claims about silent structure: LF recycling

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Chung, Ladusaw and McCloskey 1995:240–6:

IP **recycling** can be thought of as copying the LF of some discourse-available IP into the empty IP position. ... In [some cases], the recycled IP does not come supplied with a syntactic position for the displaced [WH] constituent to bind. When such a position does not already exist, it must be created, by an additional part of the recycling process we call **sprouting**.

- Case matching: OK: The WH is base-generated, and must bind (be coindexed with) some DP inside the reconstructed sluice. This kind of binding must be sensitive to case.
- Island insensitivity: 😊 Being bound is not island-sensitive.
- Sprouting: Well... As long as the reconstructed LF obeys all of the selectional and other syntactic constraints of antecedent, sprouting is ok (see quotation above).

## Claims about silent structure: PF Deletion

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Merchant 2001 (PF Deletion): Sluicing involves movement of a wh-phrase out of a sentential [IP or FocP] constituent... followed by deletion of that node.

**Mutual entailment restriction:** clause can be deleted iff the antecedent and the deletion target mutually entail each other, modulo existential focus-closure.

- Case matching: 😊 Since the WH originated in-situ, then moved, it will show all of the case matching properties of ordinary wh-movement.
- Island insensitivity: Well... Must decide that remaining unpronounced rescues island violations
- Sprouting: 😊 There is no such thing as sprouting distinct from other types of sluicing. Generate any sluice you want; as long as it mutually entails the existential focus closure of the antecedent, no problem.

Voice alternations...

## Jäger's 2001, 2005 anaphoric approach

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$$\frac{X \Rightarrow M : A \quad Y, x : A, Z, y : B, W \Rightarrow N : C}{Y, X, Z, w : B | A, W \Rightarrow N[M/x][wM/y] : C} \quad [[L]]$$

$$\frac{X, x : A, Y \Rightarrow M : B}{X, y : A | C, Y \Rightarrow \lambda z.M[yz/x] : B | C} \quad [[R]]$$

$$\frac{X, x : A, Y \Rightarrow M : B}{X, y : C \rightsquigarrow A, Y \Rightarrow \lambda z.M[yz/x] : C \rightsquigarrow B} \quad [\rightsquigarrow]$$

- A cup moved
  - $a - \lambda P x P x . x : (np \rightsquigarrow np) / n$
  - $y : (np \rightsquigarrow np) / n, z : n, w : np \setminus s \Rightarrow \lambda u.w(yzu) : np \rightsquigarrow s$ 
    - which cup moved
    - which -  $q / (s \uparrow np) / n$
- which -  $q | (np \rightsquigarrow s) / n : \lambda P Q ? x . P x \wedge Q^+ x$

## Jäger's 2001, 2005 anaphoric approach, cont'd

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- Indefinites contribute an unbound variable.
- Presence of unbound variables must be registered on category of containing clause (e.g., 'S<sup>DP</sup>').
- WH words (e.g., *who*) ambiguous between normal version and a sluice version anaphoric to S<sup>DP</sup>.

Status with respect to the three puzzles:

- Case matching: OK: Some anaphora must be sensitive to case (S<sup>DP<sub>ACC</sub></sup>).
  - Island insensitivity: 😊 unbound variables insensitive to islands.
  - Sprouting: Oops! Analysis requires overt indefinite inner antecedent.
- (8) Even overt inner antecedents need not be indefinite:  
[John or Mary] left, but I don't know which one. (AnderBois)

- Inner antecedent must take scope over the antecedent clause.
- Sluicetag silent proform anaphoric to scope remnant
- Case matching: OK: Some anaphora must be sensitive to case.
- Island insensitivity: ☺ scopability independent of syntactic islands
- Sprouting: ☺ Reasonable assumptions explain sprouting

Summary of theoretical landscape:

	Case matching	Island insensitivity	Sprouting
LF Copying	OK	☺	Well ...
PF Deletion	☺	Well ...	☺
Indef. Anaphora	OK	☺	Oops!
Anaphora to continuation	OK	☺	☺

Quantificational binding as parasitic scope

An analysis inspired by a parallel proposal in Morrill, Fadda & Valentín 2007:52:  $he = \lambda \kappa \lambda x. \kappa x x : (DP \setminus S) // (DP \setminus (DP \setminus S))$ .

$$\begin{aligned}
 & \frac{DP \cdot (\text{said} \cdot (DP \cdot \text{left})) \vdash S}{DP \circ \lambda x(x \cdot (\text{said} \cdot (DP \cdot \text{left}))) \vdash S} \equiv \\
 & \frac{\lambda x(x \cdot (\text{said} \cdot (DP \cdot \text{left}))) \vdash DP \setminus S}{DP \circ \lambda y \lambda x(x \cdot (\text{said} \cdot (y \cdot \text{left}))) \vdash DP \setminus S} \equiv \frac{DP \setminus S \vdash DP \setminus S \quad S \vdash S}{S // (DP \setminus S) \circ (DP \setminus S) \vdash S} \setminus L \\
 & \frac{\lambda y \lambda x(x \cdot (\text{said} \cdot (y \cdot \text{left}))) \vdash DP \setminus (DP \setminus S)}{\text{everyone} \circ ((DP \setminus S) // (DP \setminus (DP \setminus S))) \circ \lambda y \lambda x(x \cdot (\text{said} \cdot (y \cdot \text{left}))) \vdash S} \setminus R \text{ LEX} \\
 & \frac{\text{everyone} \circ ((DP \setminus S) // (DP \setminus (DP \setminus S))) \circ \lambda y \lambda x(x \cdot (\text{said} \cdot (y \cdot \text{left}))) \vdash S}{\text{everyone} \circ (\text{he} \circ \lambda y \lambda x(x \cdot (\text{said} \cdot (y \cdot \text{left})))) \vdash S} \setminus L \text{ LEX} \\
 & \frac{\text{everyone} \circ \lambda x(x \cdot (\text{said} \cdot (\text{he} \cdot \text{left}))) \vdash S}{\text{everyone} \cdot (\text{said} \cdot (\text{he} \cdot \text{left})) \vdash S} \equiv \\
 & \text{everyone}((\lambda \kappa \lambda x. \kappa x x)(\lambda y \lambda x. \text{said}(\text{left } x) y)) \\
 & = \text{everyone}(\lambda z. \text{said}(\text{left } z) z) = (\lambda P \forall x. P x)(\lambda z. \text{said}(\text{left } z) z) \\
 & = \forall x. \text{said}(\text{left } x) x
 \end{aligned}$$

Verb phrase ellipsis (VPE) as parasitic scope

DP<sub>he</sub>:  $\lambda \kappa \lambda x. \kappa x x : (DP \setminus S) // (DP \setminus (DP \setminus S))$   
 VPE:  $\lambda \kappa \lambda x. \kappa x x : ((DP \setminus S) \setminus S) // ((DP \setminus S) \setminus ((DP \setminus S) \setminus S))$

- (13) a. John left or Bill did. **Basic VPE**  
 $\frac{\text{left} \circ (\text{VPE} \circ \lambda y \lambda x((\text{John} \cdot x) \cdot (\text{or} \cdot (\text{Bill} \cdot y)))) \vdash S}{\text{left} \circ \lambda x((\text{John} \cdot x) \cdot (\text{or} \cdot (\text{Bill} \cdot \text{VPE}))) \vdash S} \equiv$   
 b.  $\frac{\text{left} \circ \lambda x((\text{John} \cdot x) \cdot (\text{or} \cdot (\text{Bill} \cdot \text{VPE}))) \vdash S}{(\text{John} \cdot \text{left}) \cdot (\text{or} \cdot (\text{Bill} \cdot \text{VPE})) \vdash S} \equiv$
- (14) a. John said he left or Bill did. **Sloppy coreference**  
 $\frac{DP \circ (\text{he} \circ \lambda y \lambda x(x \cdot (\text{said} \cdot (y \cdot \text{left})))) \vdash S}{DP \circ \lambda x(x \cdot (\text{said} \cdot (\text{he} \cdot \text{left}))) \vdash S} \equiv$   
 b.  $\frac{DP \circ \lambda x(x \cdot (\text{said} \cdot (\text{he} \cdot \text{left}))) \vdash S}{DP \cdot (\text{said} \cdot (\text{he} \cdot \text{left})) \vdash S} \equiv \setminus R$   
 c. Use this VP in place of *left* in (13); semantic value  $\lambda x. \text{said}(\text{left } x) x$
- (15) a. John said he left or Bill did. **Strict coreference**  
 $\frac{\text{John} \circ (\text{he} \circ \lambda y \lambda x((x \cdot (\text{said} \cdot (y \cdot \text{left}))) (\text{or} \cdot (\text{Bill} \cdot \text{VPE})))) \vdash S}{\text{John} \circ \lambda x((x \cdot (\text{said} \cdot (\text{he} \cdot \text{left}))) (\text{or} \cdot (\text{Bill} \cdot \text{VPE}))) \vdash S} \equiv$   
 b.  $\frac{\text{John} \circ \lambda x((x \cdot (\text{said} \cdot (\text{he} \cdot \text{left}))) (\text{or} \cdot (\text{Bill} \cdot \text{VPE}))) \vdash S}{(\text{John} \cdot (\text{said} \cdot (\text{he} \cdot \text{left}))) (\text{or} \cdot (\text{Bill} \cdot \text{VPE})) \vdash S} \equiv$   
 c. Continue the proof by using the VPsaid *y left* to bind VPE.

Basic sluicing

SLUICETAG:  $\lambda \kappa \lambda P. \kappa P P : ((DP \setminus S) \setminus S) // ((DP \setminus S) \setminus ((DP \setminus S) \setminus S))$

(16) Someone left, but I don't know who SLUICETAG.

The continuation of *someone* relative to the clause *someone left* (i.e.,  $\lambda x(x \cdot \text{left})$ ) provides the semantic value for the sluice gap:

$$\begin{aligned}
 & \frac{(\text{someone} \circ DP \setminus S) \cdot (\text{bidk} \cdot (\text{who} \cdot DP \setminus S)) \vdash S}{DP \setminus S \circ \lambda y((\text{someone} \circ y) \cdot (\text{bidk} \cdot (\text{who} \cdot DP \setminus S))) \vdash S} \equiv \\
 & \frac{\lambda y((\text{someone} \circ y) \cdot (\text{bidk} \cdot (\text{who} \cdot DP \setminus S))) \vdash (DP \setminus S) \setminus S}{DP \setminus S \circ \lambda z \lambda y((\text{someone} \circ y) \cdot (\text{bidk} \cdot (\text{who} \cdot z))) \vdash (DP \setminus S) \setminus S} \equiv \frac{DP \cdot DP \setminus S \vdash S}{DP \circ \lambda x(x \cdot \text{left}) \vdash S} \equiv \\
 & \frac{\lambda z \lambda y((\text{someone} \circ y) \cdot (\text{bidk} \cdot (\text{who} \cdot z))) \vdash (DP \setminus S) \setminus S}{\lambda z \lambda y((\text{someone} \circ y) \cdot (\text{bidk} \cdot (\text{who} \cdot z))) \vdash (DP \setminus S) \setminus ((DP \setminus S) \setminus S)} \setminus R \equiv \frac{\lambda x(x \cdot \text{left}) \vdash DP \setminus S \quad S \vdash S}{\lambda x(x \cdot \text{left}) \circ (DP \setminus S) \setminus S \vdash S} \setminus L \\
 & \frac{\lambda x(x \cdot \text{left}) \circ (((DP \setminus S) \setminus S) // ((DP \setminus S) \setminus ((DP \setminus S) \setminus S))) \circ \lambda z \lambda y((\text{someone} \circ y) \cdot (\text{bidk} \cdot (\text{who} \cdot z))) \vdash S}{\lambda x(x \cdot \text{left}) \circ (\text{SLUICETAG} \circ \lambda z \lambda y((\text{someone} \circ y) \cdot (\text{bidk} \cdot (\text{who} \cdot z)))) \vdash S} \setminus L \text{ LEX} \\
 & \frac{\lambda x(x \cdot \text{left}) \circ \lambda y((\text{someone} \circ y) \cdot (\text{bidk} \cdot (\text{who} \cdot \text{SLUICETAG}))) \vdash S}{(\text{someone} \circ \lambda x(x \cdot \text{left})) \cdot (\text{bidk} \cdot (\text{who} \cdot \text{SLUICETAG})) \vdash S} \equiv \\
 & \frac{(\text{someone} \cdot \text{left}) \cdot (\text{bidk} \cdot (\text{who} \cdot \text{SLUICETAG})) \vdash S}{(\text{someone} \cdot \text{left}) \cdot (\text{bidk} \cdot (\text{who} \cdot \text{SLUICETAG})) \vdash S} \equiv
 \end{aligned}$$

*bidk* = but-I-don't-know

## Good prediction: scope of inner antecedent

29/42

CLM p. 255 [my paraphrase]:

Inner antecedents must take scope over the rest of the antecedent.

(17) Each student wrote a paper on a Mayan language,  
but I don't remember which one. [CLM]

(18) Someone saw everyone, but I don't know who.

(16) Ann photographed a woman and/\*or a building yesterday, but I  
don't know who

(17) \*No one spoke to a neighbor of his, but I don't know who.

(18) Every teacher called more than two students. [\*more-than-two >  
every]

(19) Every teacher called more than two students, but I don't know who.

30/42

## Good prediction: no syntactic island sensitivity

- The relationship between the inner antecedent and the antecedent clause is scopability, not wh-extractability.
- Indefinites in particular can scope out of syntactic islands.

## Case matching

31/42

(19) who:  $Q/(DP_{ACC} \setminus S)$   
 $Q/(DP_{DAT} \setminus S)$

(20) a. SLUICEGAP:  $((DP_{ACC} \setminus S) \setminus S) // ((DP_{ACC} \setminus S) \setminus ((DP_{ACC} \setminus S) \setminus S))$   
b.  $((DP_{DAT} \setminus S) \setminus S) // ((DP_{DAT} \setminus S) \setminus ((DP_{DAT} \setminus S) \setminus S))$   
c. pn:  $(DP_F \setminus S) // (DP_F \setminus (DP_F \setminus S))$   
d.  $(DP_M \setminus S) // (DP_M \setminus (DP_M \setminus S))$

As in Jäger 2001, given an anaphoric type-logical treatment,

“Sluicing is correctly predicted to be insensitive to syntactic islands, but sensitive to morphological features of the antecedent.”

**Full accounting principle of category formation:** As in Jacobson (e.g., 1999), the category of a larger expression registers information about its missing pieces: there is no hiding of information in the derivational history.

## Sprouting: a simple case

32/42

Suggested independently to me by Lucas Champollion and Dylan Bumford: If (some) WH phrases were S modifiers (rather than VP modifiers), the analysis would extend to sprouting immediately.

(21) a. I want to know why John left.  
b. I want to know why Mary said John left. (unambiguous)  
c. *why*:  $S/S$ ; WHYSLUICEGAP:  $(S \setminus S) // (S \setminus (S \setminus S))$   
d. Target: Mary said John left, but I don't know why.

$$\frac{(\text{John} \cdot \text{left}) \circ (\text{WHYSLUICEGAP} \circ \lambda y \lambda x ((\text{Mary} \cdot (\text{said} \cdot x)) \cdot (\text{bidk} \cdot (\text{why} \cdot y)))) \vdash S}{(\text{John} \cdot \text{left}) \circ \lambda x ((\text{Mary} \cdot (\text{said} \cdot x)) \cdot (\text{bidk} \cdot (\text{why} \cdot \text{WHYSLUICEGAP}))) \vdash S} \equiv$$

$$\frac{\quad}{(\text{Mary} \cdot (\text{said} \cdot (\text{John} \cdot \text{left}))) \cdot (\text{bidk} \cdot (\text{why} \cdot \text{WHYSLUICEGAP})) \vdash S} \equiv$$

For the other reading, take *Mary said John left* as the antecedent.

Perfectly straightforward anaphora to a clause.



- (22) a. I want to know when Mary said John left. (ambiguous!)  
 b. *when*:  $S / (ADV \backslash S)$ , where  $ADV = (DP \backslash S) \backslash (DP \backslash S)$   
 c.  $WHENSLGAP: ((ADV \backslash S) \backslash S) // ((ADV \backslash S) \backslash ((ADV \backslash S) \backslash S))$   
 d. Target: Mary said John left, but I don't know when [she said he (left -)].  
 e. Need to find an ADV position inside of *John left*.

- Strategy: allow empty antecedents
- Empty antecedents usually avoided in TLG (\*very man)
- Silent lexical entries avoided in general
- Strategies for eliminating silence, as in Jäger, could be tried;
- ...if so, however, unsure about interaction with swiping.
- In any case, already using silent lexical entry for SLUICEGAP.

$$\begin{aligned}
 & \text{when} : Q / (ADV \backslash S), \text{WHENSLGAP} = ((ADV \backslash S) \backslash S) // ((ADV \backslash S) \backslash ((ADV \backslash S) \backslash S)) \\
 & \text{ADV} = (DP \backslash S) \backslash (DP \backslash S) \\
 & \frac{(DP \backslash S) \vdash DP \backslash S}{(DP \backslash S) \cdot () \vdash DP \backslash S} \equiv \\
 & \frac{}{() \vdash (DP \backslash S) \backslash (DP \backslash S)} \backslash R \\
 & \frac{}{() \vdash ADV} \text{DEF} \quad \frac{S \cdot (\text{bidk} \cdot (\text{when} \cdot ADV \backslash S)) \vdash S}{((() \circ ADV \backslash S) (\text{bidk} \cdot (\text{when} \cdot ADV \backslash S))) \vdash S} \backslash L \\
 & \frac{ADV \backslash S \circ \lambda y ((() \circ y) (\text{bidk} \cdot (\text{when} \cdot ADV \backslash S))) \vdash S}{\lambda y ((() \circ y) (\text{bidk} \cdot (\text{when} \cdot ADV \backslash S))) \vdash (ADV \backslash S) \backslash S} \backslash L \\
 & \frac{\lambda y ((() \circ y) (\text{bidk} \cdot (\text{when} \cdot z))) \vdash (ADV \backslash S) \backslash S}{ADV \backslash S \circ \lambda z \lambda y ((() \circ y) (\text{bidk} \cdot (\text{when} \cdot z))) \vdash (ADV \backslash S) \backslash S} \backslash R \\
 & \frac{\lambda z \lambda y ((() \circ y) (\text{bidk} \cdot (\text{when} \cdot z))) \vdash (ADV \backslash S) \backslash ((ADV \backslash S) \backslash S)}{\lambda x (\text{John} \cdot (\text{left} \cdot x)) \circ (\text{WHENSLGAP} \circ \lambda z \lambda y ((() \circ y) (\text{bidk} \cdot (\text{when} \cdot z)))) \vdash S} \backslash L \\
 & \frac{\lambda x (\text{John} \cdot (\text{left} \cdot x)) \circ \lambda y ((() \circ y) (\text{bidk} \cdot (\text{when} \cdot \text{WHENSLGAP}))) \vdash S}{((() \circ \lambda x (\text{John} \cdot (\text{left} \cdot x))) \cdot (\text{bidk} \cdot (\text{when} \cdot \text{WHENSLGAP}))) \vdash S} \equiv \\
 & \frac{(\text{John} \cdot (\text{left} \cdot ())) \cdot (\text{bidk} \cdot (\text{when} \cdot \text{WHENSLGAP})) \vdash S}{(\text{John} \cdot \text{left}) \cdot (\text{bidk} \cdot (\text{when} \cdot \text{WHENSLGAP})) \vdash S} \equiv
 \end{aligned}$$

Independent motivation for empty antecedents: deriving gaps <sup>34/42</sup>

- Assume the empty structure, '()', is an identity element for  $\circ$
- So  $\Gamma \circ () \equiv \Gamma \equiv () \circ \Gamma$

$$\frac{DP \backslash S \vdash DP \backslash S}{() \circ DP \backslash S \vdash DP \backslash S} \equiv \\
 \frac{}{() \vdash (DP \backslash S) // (DP \backslash S)} // R$$

who:  $Q / (DP \backslash S)$ : who does John like:

$$\frac{\text{does} \cdot (\text{John} \cdot (\text{like} \cdot DP)) \vdash S}{DP \circ \lambda x (\text{does} \cdot (\text{John} \cdot (\text{like} \cdot x))) \vdash S} \equiv \\
 \frac{\lambda x (\text{does} \cdot (\text{John} \cdot (\text{like} \cdot x))) \vdash DP \backslash S}{(DP \backslash S) // (DP \backslash S) \circ \lambda x (\text{does} \cdot (\text{John} \cdot (\text{like} \cdot x))) \vdash DP \backslash S} \backslash R \\
 \frac{}{\text{GAP} \circ \lambda x (\text{does} \cdot (\text{John} \cdot (\text{like} \cdot x))) \vdash DP \backslash S} \text{LEX} \\
 \frac{}{\text{does} \cdot (\text{John} \cdot (\text{like} \cdot \text{GAP})) \vdash DP \backslash S} \equiv$$

Likewise for  $\cdot$  mode. Silent elements usually avoided in TLG, but standard in many logical settings.

Implicit arguments

- (23) a. John ate, but I don't know what.  
 b. New category: given  $A, B$  formulas,  $A \otimes B$   
 c. Residuation laws:  $A \vdash C / B$  iff  $A \otimes B \vdash C$  iff  $B \vdash A \backslash C$   
 d.  $\text{ate}_{\text{INTR}} : \langle \text{eat}_{\text{tr}}, \lambda P \exists x. Px \rangle : ((DP \backslash S) / DP) \otimes S // (DP \backslash S)$

$$\frac{\Sigma [A \cdot B] \vdash C}{\Sigma [A \otimes B] \vdash C} \otimes L \quad \frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma \cdot \Delta \vdash A \otimes B} \otimes R$$

$$\frac{(\text{John} \cdot (((DP \backslash S) / DP) \cdot S // (DP \backslash S))) \cdot (\text{bidk} \cdot (\text{what} \cdot \text{SLUICEGAP})) \vdash S}{(\text{John} \cdot ((DP \backslash S) / DP) \otimes S // (DP \backslash S)) \cdot (\text{bidk} \cdot (\text{what} \cdot \text{SLUICEGAP})) \vdash S} \otimes L \\
 \frac{}{(\text{John} \cdot \text{ate}_{\text{INTRANS}}) \cdot (\text{bidk} \cdot (\text{what} \cdot \text{SLUICEGAP})) \vdash S} \text{LEX}$$

- (24) a. Everyone ate, but I don't know what.  $\forall > \exists, ?* \exists > \forall$   
 b. ?No one ate, but I don't know what.

Available to Jäger; how to guarantee narrowest scope of IA?

## Problems for mutual entailment

37/42

Romero, Merchant: the focus closure of the antecedent clause and the sluice must entail each other.

Counterexamples:

- (20) \*Kelly was murdered, but we don't know who.
- (21) \*Someone paid Mary, but we don't know by whom.
- (22) Some numbers between 2 and 20 are even or odd, but I'm not going to tell you which numbers are ~~prime or not prime~~.

## The answer ban

39/42

- The antecedent clause must not resolve (or partly resolve) the issue raised by the sluiced interrogative.

- (27) \*John left, but I don't know who.
- (28) John left, but I don't know who else.
- (29) \*John or Mary left, but I don't know who.
- (30) John met a woman, but I don't know who.
- (31) Mary knows that John left, but Bill doesn't know who.

## The wh-correlate does NOT need to be indefinite

38/42

- (23) I know that John left, but I don't know who else.
- (24) Mary has dined at Masa, and I don't know where else.
- (25) John liked the collards, but I don't know which other dishes.
- (26) Mary tasted each hot dish, and I don't know what else.

## Andrews Amalgams: ellipsis to a containing continuation<sup>40/42</sup>

- (33) Johnson 2013:
  - a. Sally will eat something today, but I don't know what ...
  - b. Sally will eat [I don't know what ...] today.

$$\begin{array}{c}
 \frac{idk \cdot (\text{what} \cdot DP \setminus S) \vdash S}{DP \setminus S \circ \lambda x(idk \cdot (\text{what} \cdot x)) \vdash S} \equiv \\
 \frac{\lambda x(idk \cdot (\text{what} \cdot x)) \vdash (DP \setminus S) \setminus S \quad G \vdash G}{G \setminus ((DP \setminus S) \setminus S) \circ \lambda x(idk \cdot (\text{what} \cdot x)) \vdash G} \setminus R \quad //L \\
 \frac{AMALGAM \circ \lambda x(idk \cdot (\text{what} \cdot x)) \vdash G}{idk \cdot (\text{what} \cdot AMALGAM) \vdash G} \equiv \\
 \frac{\lambda y(idk \cdot (\text{what} \cdot y)) \vdash (DP \setminus S) \setminus S \quad G \circ \lambda x(Sally \cdot (\text{ate} \cdot x)) \vdash S}{(G \setminus ((DP \setminus S) \setminus S) \circ \lambda y(idk \cdot (\text{what} \cdot y))) \circ \lambda x(Sally \cdot (\text{ate} \cdot x)) \vdash S} \setminus L \\
 \frac{(idk \cdot (\text{what} \cdot AMALGAM)) \circ \lambda x(Sally \cdot (\text{ate} \cdot x)) \vdash S}{Sally \cdot (\text{ate} \cdot (idk \cdot (\text{what} \cdot AMALGAM))) \vdash S} \equiv, \text{LEX}
 \end{array}$$

$G \equiv S \setminus (DP \setminus S)$  (i.e., scope-taking DP, a generalized quantifier)

## Mismatching examples

41/42

Chung 2006: The syntactic objects which are copied or re-used will have to be abstract enough to permit certain ‘mismatches’ between the antecedent and the apparent requirements of the ellipsis-site.

- (25) a. John remembers meeting someone,  
but he doesn't remember who [~~he met~~].  
b.  $((DP \setminus S_{-ING}) \setminus S) // ((DP \setminus S) \setminus ((DP \setminus S_{-ING}) \setminus S))$

- Syntax is no problem.
- Semantically, no need to build a tensed clause: only necessary to turn an -ING clause meaning into a tensed clause meaning.
- In this case, we need a function from a “remembering” event type to an open proposition concerning a specific event within that event type

## Claims

42/42

- The ellipsis site contains a silent proform, e.g., SLUICEGAP
- So silent elements are ok—but don't have internal structure
- The syntactic category of the inner antecedent is transparently available to the sluicegap, **case matching is easy**
- The inner antecedent must scope over the antecedent clause
- Because the only constraint on the relationship between the inner antecedent and the antecedent clause is scopability, sluicing is **insensitive to syntactic islands**.
- When implemented by a suitable type logical grammar that allows reasoning about scope, **sprouting follows** from independently motivated assumptions about empty antecedents

**Sluicing is anaphora to an anti-constituent, that is, anaphora to a continuation.**