# Syntactic tiers for movement and agreement 

## Day 3: Tree Rewriting \& Externalization

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## Outline

1 Subcategorization is too powerful

2 Feature recoverability as strictly local rewriting

3 Bare phrase structure: Local transductions for Merge and Move

4 The challenge of linearization

## Take-home message

- Overgeneration problem in syntax

Subcategorization can express very unnatural constraints, due to category refinement.

- A linguistically fertile solution

Category features don't come for free. They must be inferable from the local context.

## Hidden power of subcategorization

Every formalism with subcategorization can express undesirable constraints. (Graf 2017)

Counting every DP contains at least five LIs
Symmetry closure every reflexive c-commands its antecedent Complement sentence well-formed jiff ill-formed in English
Boolean closure sentence must obey either V2 or Principle A, unless there are less than 7 pronounced LIs
Domain blindness a sentence is well-formed of there are at least two words that display word-final devoicing
Is(n't)lands an adjunct is an island jiff
it is inside an embedded clause or
it contains no animate nouns

## Why?

- Complex constraints can be lexicalized by decomposing them into refined categories.
- They are then enforced via subcategorization.
- It's a generalized version of slash feature percolation. (Gazdar et al. 1985; Graf 2011; Kobele 2011)


## An example

## Subcategorization (Stabler 1997)

- Category features ( $\mathrm{F}^{-}$)
- Selector features ( $\mathrm{F}^{+}$)
- Subcategorization: matching features of opposite polarity

A very simple grammar

$$
\begin{array}{rr}
\text { foo }:: \mathrm{X}^{-} & \text {foo }:: \mathrm{X}^{+} \mathrm{X}^{-} \\
\text {bar }:: \mathrm{X}^{-} & \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} \\
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-}
\end{array}
$$

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\text {bar }:: \mathrm{X}^{-} \quad \text { bar }:: \mathrm{X}^{+} \mathrm{X}^{-} \\
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-} \quad \text { foo }:: \mathrm{X}^{-}
\end{array}
$$

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\text {bar }:: \mathrm{X}^{-} & \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {। } \\
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-} & \text {foo }:: \mathrm{X}^{-}
\end{array}
$$

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- Category features ( $\mathrm{F}^{-}$)
- Selector features ( $\mathrm{F}^{+}$)
- Subcategorization: matching features of opposite polarity

A very simple grammar

$$
\begin{array}{ccc} 
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-} \\
\text {foo :: } \mathrm{X}^{-} & \text {foo }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {bar }:: \\
\text { bar }:: \mathrm{X}^{+} \mathrm{X}^{-} \\
& \text {bar :: } \mathrm{X}^{+} \mathrm{X}^{-} & \quad 1 \\
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-} & \text {foo :: } \mathrm{X}^{-}
\end{array}
$$

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\text {bar :: } \mathrm{X}^{-} & \text {bar :: } \mathrm{X}^{+} \mathrm{X}^{-} & \\
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-} & \text {foo :: } \mathrm{X}^{-} \\
& & \\
& &
\end{array}
$$

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\begin{array}{rrrr} 
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\text {foo }:: \mathrm{X}^{-} & \text {foo }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} & \\
\text {bar }:: \mathrm{X}^{-} \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {। } & \\
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-} & \text {foo }:: \mathrm{X}^{-} & \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} \\
& & & \text {। } \\
& & & \text { foo }:: \mathrm{X}^{-}
\end{array}
$$

## An example

## Subcategorization (Stabler 1997)

- Category features ( $\mathrm{F}^{-}$)
- Selector features $\left(\mathrm{F}^{+}\right)$
- Subcategorization: matching features of opposite polarity

A very simple grammar

$$
\begin{array}{rrrr} 
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-} \\
\text {foo }:: \mathrm{X}^{-} & \text {foo }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} \\
\text {bar }:: \mathrm{X}^{-} & \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {। } & \text { । } \\
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-} & \text {foo }:: \mathrm{X}^{-} & \text {bar :: } \mathrm{X}^{+} \mathrm{X}^{-} \\
& & & \text {। } \\
& & & \text { foo }:: \mathrm{X}^{-}
\end{array}
$$

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\text {foo }:: \mathrm{X}^{-} & \text {foo }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {bar :: } \mathrm{X}^{+} \mathrm{X}^{-} \\
\text {bar }:: \mathrm{X}^{-} & \text {bar }:: \mathrm{X}^{+} \mathrm{X}^{-} & \text {। } & \text { । } \\
& \varepsilon:: \mathrm{X}^{+} \mathrm{C}^{-} & \text {foo }:: \mathrm{X}^{-} & \text {bar :: } \mathrm{X}^{+} \mathrm{X}^{-} \\
& & & \text {foo :: } \mathrm{X}^{-}
\end{array}
$$

## Adding modulo counting

- Suppose every tree must have an even number of nodes
- Refinement: $\mathrm{X}^{-} \Rightarrow \mathrm{O}^{-}$and $\mathrm{E}^{-}$for Odd and Even

Refined grammar with even/odd distinction

$$
\begin{array}{rr}
\text { foo }:: \mathrm{O}^{-} & \text {foo }:: \mathrm{E}^{+} \mathrm{O}^{-} \\
& \text {foo }:: \mathrm{O}^{+} \mathrm{E}^{-} \\
\text {bar }:: \mathrm{O}^{-} & \text {bar }:: \mathrm{E}^{+} \mathrm{O}^{-} \\
& \text {bar }:: \mathrm{O}^{+} \mathrm{E}^{-} \\
& \varepsilon:: \mathrm{O}^{+} \mathrm{C}^{-}
\end{array}
$$

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- Suppose every tree must have an even number of nodes
- Refinement: $\mathrm{X}^{-} \Rightarrow \mathrm{O}^{-}$and $\mathrm{E}^{-}$for Odd and Even

Refined grammar with even/odd distinction

$$
\begin{aligned}
& \text { foo :: } 0^{-} \text {foo :: } \mathrm{E}^{+} \mathrm{O}^{-} \\
& \text {foo :: } 0^{+} \mathrm{E}^{-} \\
& \text {bar :: } \mathrm{O}^{-} \text {bar :: } \mathrm{E}^{+} \mathrm{O}^{-} \\
& \text {bar :: } \mathrm{O}^{+} \mathrm{E}^{-} \text {foo :: } \mathrm{O}^{-} \\
& \varepsilon:: 0^{+} \mathrm{C}^{-}
\end{aligned}
$$

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\begin{array}{rrrr}
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& \text {foo }:: \mathrm{O}^{+} \mathrm{E}^{-} & \text {bar :: } 0^{+} \mathrm{E}^{-} \\
\text {bar :: } \mathrm{O}^{-} & \text {bar :: } & \mathrm{E}^{+} \mathrm{O}^{-} & \text {। } \\
& \text { bar }:: & \mathrm{O}^{+} \mathrm{E}^{-} & \text {foo }:: 0^{-} \\
& \varepsilon:: & \mathrm{O}^{+} \mathrm{C}^{-} &
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& \text {foo }:: \mathrm{O}^{+} \mathrm{E}^{-} & \text {bar }:: \mathrm{O}^{+} \mathrm{E}^{-} \\
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& \text { bar }:: \mathrm{O}^{+} \mathrm{E}^{-} & \text {foo }:: \mathrm{O}^{-} \\
& \varepsilon: & \mathrm{O}^{+} \mathrm{C}^{-}
\end{array}
$$

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& \text {foo }:: \mathrm{O}^{+} \mathrm{E}^{-} & \text {bar }:: \mathrm{O}^{+} \mathrm{E}^{-} \\
\text {bar }:: \mathrm{O}^{-} & \text {bar }:: \mathrm{E}^{+} \mathrm{O}^{-} & \text {। } \\
& \text { bar }:: \mathrm{O}^{+} \mathrm{E}^{-} & \text {foo }:: \mathrm{O}^{-} \\
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\text {bar }:: \mathrm{O}^{-} & \text {bar }:: \mathrm{E}^{+} \mathrm{O}^{-} & \text {। } & \\
& \text { bar }:: \mathrm{O}^{+} \mathrm{E}^{-} & \text {foo }:: \mathrm{O}^{-} & \text {bar }:: \mathrm{O}^{+} \mathrm{E}^{-} \\
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& \text {foo }:: \mathrm{O}^{+} \mathrm{E}^{-} & \text {bar }:: \mathrm{O}^{+} \mathrm{E}^{-} & \text {bar }:: \mathrm{E}^{+} \mathrm{O}^{-} \\
\text {bar }:: \mathrm{O}^{-} & \text {bar }:: \mathrm{E}^{+} \mathrm{O}^{-} & \text {। } & \text { foo }:: \mathrm{O}^{-} & \text {bar }:: \mathrm{O}^{+} \mathrm{E}^{-} \\
& \text {bar }:: \mathrm{O}^{+} \mathrm{E}^{-} & & \text {foo }:: \mathrm{O}^{-}
\end{array}
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& \text {bar }:: \mathrm{O}^{+} \mathrm{E}^{-} \\
& \varepsilon:: \mathrm{O}^{+} \mathrm{C}^{-}
\end{array}
$$

## The problem with subcategorization

- Even very complex constraints can be

1 compiled into the category system and
2 enforced via subcategorization.

- works for all MSO constraints $\Rightarrow$ massive overgeneration (Graf 2011; Kobele 2011)
- Linguistic criteria for determining categories are too weak to prevent this.
- morphology
- syntactic distribution
- semantics


## The central issue

We need a more restrictive notion of category!

## A formal notion of complexity

- We need to restrict the power of subcategorization, but how?
- Features currently come for free.
- We must measure the cost of features.


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## A formal notion of complexity

- We need to restrict the power of subcategorization, but how?
- Features currently come for free.
- We must measure the cost of features.
water :: $\mathrm{D}^{+} \mathrm{D}^{+} \mathrm{V}^{-} \longleftarrow$ feature assignment water


Local feature recoverability
Features must be recoverable in an ISL fashion.

## Input strictly k-local relabelings

## ISL string-to-string transduction (Chandlee 2014)

Rewrite each symbol in a string based on its local input context.

## An ISL-3 relabeling


b b a
a c a e

- ISL is a very weak class, yet widely found in phonology and morphology.


## Input strictly k-local relabelings

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Rewrite each symbol in a string based on its local input context.

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b b
a
a c a e

$$
a \quad b \quad b \quad c \quad a
$$

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$\square$

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$a \quad b \quad b \quad a \quad a$
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a \quad b \quad b \quad a \quad a
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## An ISL-3 relabeling



$$
a \quad b \quad b \quad a \quad a
$$

$$
b \quad b \quad a \quad b
$$

$$
a \quad c \quad a \quad e
$$

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## Input strictly k-local relabelings

## ISL string-to-string transduction (Chandlee 2014)

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$$
a \quad b \quad b \quad a \quad a
$$

b b a b e
a c a e

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## Input strictly k-local relabelings

## ISL string-to-string transduction (Chandlee 2014)

Rewrite each symbol in a string based on its local input context.

## An ISL-3 relabeling



$$
\begin{array}{llllll}
a & b & b & a & c & a \\
b & b & a & b & e & b
\end{array}
$$

$$
a \quad c \quad a \quad e
$$

- ISL is a very weak class, yet widely found in phonology and morphology.


## Lifting ISL relabelings to trees

## String contexts as tree contexts

a cae

## Lifting ISL relabelings to trees

## String contexts as tree contexts

a $c$ a e
$a-c=a$
e

## Lifting ISL relabelings to trees

## String contexts as tree contexts



## Lifting ISL relabelings to trees

## String contexts as tree contexts

a $C$ a e
$a-c-a$ e



## Reminder: ISL for feature inference

- Feature cost $\approx$ how hard to assign by transduction?

$$
\text { water }:: \mathrm{D}^{+} \mathrm{D}^{+} \mathrm{V}^{-} \longleftarrow \text { feature assignment } \text { water }
$$



Local feature recoverability
Features must be recoverable in an ISL fashion.

## Reminder: ISL for feature inference

- Feature cost $\approx$ how hard to assign by transduction?

$$
\text { water :: } \mathrm{D}^{+} \mathrm{D}^{+} \mathrm{V}^{-}
$$



## Local feature recoverability

Features must be recoverable in an ISL fashion.

## Intuition

Categorial ambiguity can be resolved within local context

## Modulo counting is not ISL recoverable



## Modulo counting is not ISL recoverable



- Can you determine the features of foo?
$10^{+} E^{-}$
$2 \mathrm{E}^{+} \mathrm{O}^{-}$
- No, that's impossible.
- You need more than local information.
- Modulo counting is not ISL recoverable.
bar


## An empirical conjecture

## SL-2 recoverability conjecture

The category and selector features of lexical items are

- recoverable from feature-less dependency trees
- using only a window of size 2 .

gardeners flowers



## Implications and open issues

## Implications

- We avoid tons of overgeneration.
- Heads only select for arguments, not arguments of arguments.

Open issues

- Needs to be tested across many languages
- Depends on theoretical assumptions
- distribution of empty heads
- lexical items fully inflected or bare roots? (Hale and Keyser 1993; Marantz 1997)
- SL-2 may be too tight, but SL- $k$ recoverability seems safe
- Move features are not ISL recoverable!


## An incomplete picture

Movement isn't just a syntactic dependency, it affects the output.


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Movement isn't just a syntactic dependency, it affects the output.


## Phrase structure trees without movement

no movement $\Rightarrow$ easy translation to phrase structure trees


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no movement $\Rightarrow$ easy translation to phrase structure trees


## Adding tier relations

- ISL limited to local contexts $\Leftrightarrow$ unbounded movement
- But: movement is local over tiers
- suffices to enrich ISL rewrite rules with tier-daughter relation

ISL rewrite rule with nom-daughter relation


## Example with subject movement



## Example with subject movement



## Example with subject movement



## Example with subject movement



## Example with subject movement



## Example with subject movement



## Example with subject movement



## Example with subject movement




## Multiple copies are straightforward



## Multiple copies are straightforward



## Multiple copies are straightforward



## Multiple copies are straightforward



## Multiple copies are straightforward



## Multiple copies are straightforward



## Multiple copies are straightforward



## Multiple copies are straightforward



## Multiple copies are straightforward



## Multiple copies are straightforward



## Linearization/traces: tricky

- Copies don't tell us how to pronounce the tree.
- Traces: unpronounced landing sites of movers



## Base delinking example


what

## Lexical predictability

- Given two landing sites $x$ and $y$ on different tiers, one cannot tell from the tiers whether x or y is higher.
- We cannot distinguish final from non-final landing sites.


Lexical predictability requirement of delinking
Delinking works only if one knows whether to
1 insert a copy, or
2 insert a trace.
Due to the limitations of tiers, this must be inferable directly from the mover.

## Empirical support

- Lexical predictability holds for nom and wh movement.


## Ban on Improper Movement (BolM)

If a mover undergoes both nom and wh movement, nom movement derivationally precedes wh movement.

might

## Empirical support

- Lexical predictability holds for nom and wh movement.


## Ban on Improper Movement (BoIM)

If a mover undergoes both nom and wh movement, nom movement derivationally precedes wh movement.

$\rightarrow$


## Output-oriented Ban on Improper Movement

- BolM is a particular instance of a more general requirement.


## Output-oriented Ban on Improper Movement (OOBoIM)

- Let I be an arbitrary lexical item with $\left\{\mathrm{f}^{-}, \mathrm{g}_{0}^{-}, \ldots, \mathrm{g}_{n}^{-}\right\}$.
- If I's final movement step is f-movement in some derivation, then I's final movement step is f-movement in all derivations.
- Kenneth Hanson's analysis of MG corpus supports even stronger version: if $f \prec g$ for I in some derivation, then $\mathrm{f} \prec \mathrm{g}$ for I in all derivations.
- OOBoIM permits BolM violations hyperraising


Kenneth Hanson

## Conclusion

- Movement is both a syntactic dependency and an operation.
- In both cases the core of movement is local over tiers.
- Identifying mover with all landing sites (copies/multidominance) easier than identifying output-relevant landing sites (traces)


## Outlook: Bringing it all together

- TSL perspectives of all movement types covert, successive-cyclic, sidewards, multiple wh
- Tiers for islands, extraction morphology and path conditions Wolof u-chains, floating quantifiers, Germanic wh-copying
- Learning

SL learning of Merge features, ??? for Move features

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